

# CS 441: Integer representation and Algorithms

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# Today's topics

- Integer representations
  - Base  $b$  expansions
  - Common bases: Binary, octal, hexadecimal
  - Base conversions
- Integer algorithms
  - Addition
  - Multiplication
  - Connection to computing and pen-and-paper arithmetic



## While we typically use decimal, other base systems work very similarly

Recall: **Decimal expansion** of integers, e.g.,

$$3528 = 3 * 10^3 + 5 * 10^2 + 2 * 10^1 + 8 * 10^0$$

**Theorem:** Let  $b$  be an integer greater than 1. Any  $n \in \mathbf{Z}^+$  can be expressed **uniquely** in the form:

$$n = a_k b^k + a_{k-1} b^{k-1} + \dots + a_1 b^1 + a_0 b^0$$

where  $k \in \mathbf{N}$ , each  $a_i \in \mathbf{N}$  where  $a_i < b$ , and  $a_k \neq 0$ .

This representation is called the **base  $b$  expansion of  $n$** , which we write compactly as  $(a_k a_{k-1} \dots a_1 a_0)_b$

- When  $b > 10$ , we write each  $a_i$  as a single symbol in an extended “alphabet” of digits
  - e.g., 0123456789ABCDEFGHIJ...

## Examples of base $b$ expansions

Express each of these expansions in decimal:

- $(675)_8 =$   
=

- $(110101)_2 =$   
=

- $(2A4)_{12} =$   
=

*Here, A has value 10*

# Common base expansions

These base systems are very common in **computing**:

- Base 2, binary: Expansions are bit strings  
 $412 = (110011100)_2$
- Base 8, octal: Each digit  $a_i$  is  $0 \leq a_i < 8$   
 $412 = (634)_8$
- Base 16, hexadecimal: Each digit  $a_i \in \{0, 1, \dots, 9, A, B, \dots, F\}$   
 $412 = (19C)_{16}$

Trends to note:

- Base  $b$  requires  $b$  digits in the “alphabet”
- Lesser  $b$  yields longer expansions, greater  $b$  yields shorter expansions

Why are these important?

- Data is stored in binary, octal represents 3 bits per digit, hexadecimal represents 4 bits per digit

# Constructing base $b$ expansions

**procedure** *base  $b$  expansion*( $n, b$ : positive integers with  $b > 1$ )

$q := n$

$k := 0$

*Digits are produced right-to-left*

**while**  $q \neq 0$

$a_k := q \bmod b$

$q := q \operatorname{div} b$

$k := k + 1$

*Repeat: Divide  $q$  by  $b$ ;  
remainder becomes a digit,  
quotient replaces  $q$*

**return**  $(a_{k-1}, a_{k-2}, \dots, a_1, a_0)$

$\{(a_{k-1}a_{k-2}\cdots a_1a_0)_b$  is the base  $b$  expansion of  $n\}$

*Return when  $q = 0$*

# Constructing base $b$ expansions, examples

## 1. Express 1501 in hex

- 1501 divided by 16  
 $q = 93, r = 13 = (D)_{16}$
- 93 divided by 16  
 $q = 5, r = 13 = (D)_{16}$
- 5 divided by 16  
 $q = 0, r = 5$
- Thus,  $1501 = (5DD)_{16}$

## 2. Express 441 in octal

- 441 divided by 8  
 $q = 55, r = 1$
- 55 divided by 8  
 $q = 6, r = 7$
- 6 divided by 8  
 $q = 0, r = 6$
- Thus,  $441 = (671)_8$

## 3. Express 441 in base-30

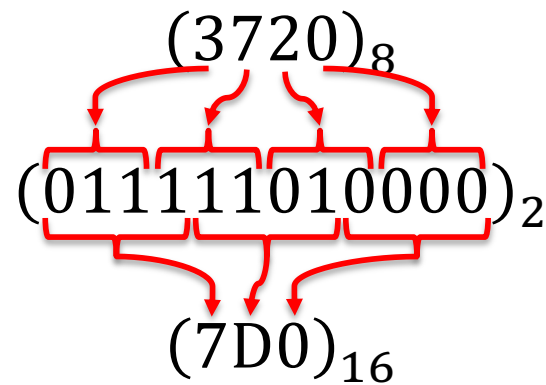
- 441 divided by 30  
 $q = 14, r = 21 = (I)_{30}$
- 14 divided by 30  
 $q = 0, r = 14 = (E)_{30}$
- Thus,  $441 = (EI)_{30}$

## 4. Express 441 in base-4

- 441 divided by 4  
 $q = 110, r = 1$
- 110 divided by 4  
 $q = 27, r = 2$
- 27 divided by 4  
 $q = 6, r = 3$
- 6 divided by 4  
 $q = 1, r = 2$
- 1 divided by 4  
 $q = 0, r = 1$
- Thus,  $441 = (12321)_4$

When  $b = 2^i$ , conversion can be done on  $i$  bits at a time

Since an octal digit encodes 3 bits and a hex digit encodes 4 bits, we can use binary to help convert





## In-class exercises

**Problem 1:** Find the octal expansion of 100

**Problem 2:** Find the octal expansion of  $(100)_2$

**Problem 3:** Find the octal expansion of  $(100)_{16}$

**Problem 4:** Find the octal expansion of  $(100)_{36}$

## Adding base $b$ expansions

```

procedure add( $x, y$ : positive integers,  $b$ : integer  $> 1$ )
  {The base  $b$  expansions of  $x$  and  $y$  are  $(x_{n-1}x_{n-2}\cdots x_1x_0)_b$ 
  and  $(y_{n-1}y_{n-2}\cdots y_1y_0)_b$ , respectively}
   $c := 0$ 
  for  $j := 0$  to  $n-1$            {Move right-to-left}
     $t := x_j + y_j + c$            {Add the  $j$ th digits together}
     $c := \lfloor t/b \rfloor$            {Carry a digit if needed}
     $s_j := t - bc$                  {Remove carry and save as  $s_j$ }
   $s_n := c$                        {Final carry becomes  $s_n$ }
  return  $(s_n, s_{n-1}, \dots, s_1, s_0)$ 
  {The base  $b$  expansion of the sum is  $(s_n s_{n-1} \cdots s_1 s_0)_b$ }

```

*Does this sound familiar?*  
*What is its complexity?*

# Addition examples in hexadecimal/octal

*Hex* | *Octal*

$$\begin{array}{r} \text{B8C0} \\ + \text{827F} \\ \hline \end{array}$$

$$\begin{array}{r} 5630 \\ + 3766 \\ \hline \end{array}$$

$$\begin{array}{r} 13AC4F \\ + 3B9E00 \\ \hline \end{array}$$

$$\begin{array}{r} 723405 \\ + 27305 \\ \hline \end{array}$$

## Multiplying base $b$ expansions

```

procedure multiply( $x, y$ : positive integers,  $b$ : integer  $> 1$ )
  {The base  $b$  expansions of  $x$  and  $y$  are  $(x_{n-1}x_{n-2}\cdots x_1x_0)_b$  and  $(y_{n-1}y_{n-2}\cdots y_1y_0)_b$ , respectively}
   $p := 0$                                      {Resulting product}
  for  $j := 0$  to  $n-1$ 
     $c := 0$                                      {Move right-to-left in  $y$ }
    for  $i := 0$  to  $n-1$ 
       $t := x_i * y_j + c$                        {Reset carry}
       $c := \lfloor t/b \rfloor$                        {Move right-to-left in  $x$ }
       $r_i := t - bc$                              {Multiply digits and add carry}
     $r_n := c$                                    {Carry a digit if needed}
     $r := r$  shifted  $j$  places                    {Partial product, digit  $i$ }
     $p := \text{add}(p, r)$                           {Final carry becomes  $r_n$ }
  return  $(p_{2n}, p_{2n-1}, \dots, p_1, p_0)$       {Shift position to align with  $j$ }
  {The base  $b$  expansion of the sum is  $(p_{2n}p_{2n-1}\cdots p_1p_0)_b$ }
  {Add  $r$  to the result}

```

**What is its complexity?**

# Multiplication examples in hexadecimal/octal

	<i>Hex</i>		<i>Octal</i>
	C38		365
	* <u>6A4</u>		* <u>457</u>
	_____		_____

# How are these algorithms used in practice?

In previous exercises, didn't we assume basic arithmetic operations were  $\Theta(1)$ ?

- This is **often** true! Modern CPUs can compute (at least) 32-bit integer multiplication in circuitry in a few cycles
- What about numbers **bigger** than your CPU's MUL?
  - e.g., for cryptography
- Let  $b = 2^{32}$ , consider  $b$ -bit expansions where each "digit" is a **32-bit word**

We can compare a CPU's MUL (etc.) circuits to the **multiplication tables** we memorized in grade school

- For small enough values, we know the answer very quickly
- For larger values, we learn an algorithm that utilizes many smaller multiplications

## Other multiplication algorithms for even larger values

Algorithm	Complexity	Threshold example*
Grade school	$O(n^2)$	(native MUL)
Karatsuba	$O(n^{\log_2 3}) \approx O(n^{1.585})$	832 bits
Toom–Cook (Toom-3)	$O(n^{\log_3 5}) \approx O(n^{1.46})$	6208 bits
Schönhage–Strassen	$O(n \log n \log \log n)$	159744 bits
Fürer	$O(n \log n 2^{\Theta(\log^* n)})$	—
Harvey–van der Hoeven	$O(n \log n)$	—

## In-class exercises

**Problem 5:** Use the integer addition algorithm to compute  $(734)_8 + (225)_8$

**Problem 6:** Use the integer multiplication algorithm to compute  $(110110)_2 \times (100101)_2$

**Problem 7:** Calculate  $(FF)_{16} \times (FF)_{16}$ ,  $(77)_8 \times (77)_8$ , and  $99 \times 99$  and compare



## Final thoughts

- Integers can be represented uniquely in any specified base
- Integer arithmetic can be computed in other bases, and even pen-and-paper algorithms can be useful in computing
  - Arithmetic isn't always constant
- Next time:
  - Primes and composites (Section 4.3)