

CS 441: Proof by induction

PhD. Nils Murrugarra-Llerena
nem177@pitt.edu



Today's topics

- Proof by Induction
- Examples
 - Equations
 - Inequalities
 - Sets



We've learned a lot of proof methods...

Basic proof methods

- Direct proof, contradiction, contraposition, cases, ...

Proof of quantified statements

- Existential statements (i.e., $\exists x P(x)$)
 - Finding a single example suffices
- Universal statements (i.e., $\forall x P(x)$) can be harder to prove

- $$\sum_{j=0}^n ar^j = \begin{cases} \frac{ar^{n+1}-a}{r-1} & \text{if } r \neq 1 \\ (n+1)a & \text{if } r = 1 \end{cases}$$

- $$\sum_{j=1}^n j = \frac{n(n+1)}{2}$$

Bottom line: We need new tools!

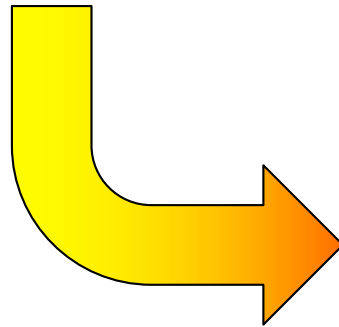
Mathematical induction lets us prove universally quantified statements!

Goal: Prove $\forall x \in \mathbf{N} P(x)$.

Intuition: If $P(0)$ is true, then $P(1)$ is true. If $P(1)$ is true, then $P(2)$ is true...

Procedure:

1. Prove $P(0)$
2. Show that $P(k) \rightarrow P(k+1)$ for any **arbitrary** k
3. Conclude that $P(x)$ is true $\forall x \in \mathbf{N}$



$P(0)$
For arb. k , $P(k) \rightarrow P(k+1)$

 $\therefore \forall x \in \mathbf{N} P(x)$

Analogy: Climbing a ladder

Proving $P(0)$:

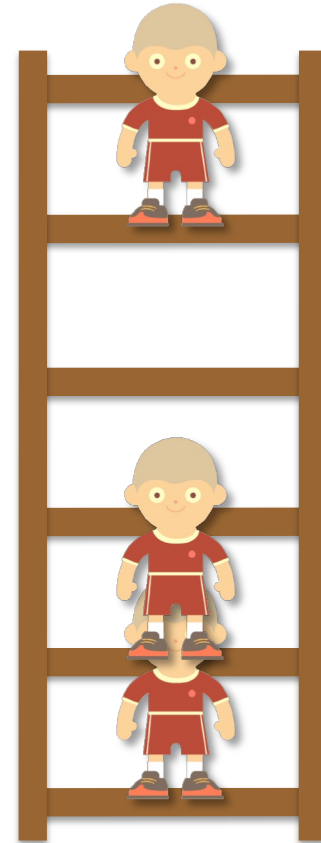
- You can get on the first rung of the ladder

Proving $P(k) \rightarrow P(k+1)$:

- If you are on the k th step, you can get to the $(k+1)$ th step

$\therefore \forall x P(x)$

- You can get to any step on the ladder



Analogy: Playing with dominoes

Proving $P(0)$:

- The first domino falls

Proving $P(k) \rightarrow P(k+1)$:

- If the k th domino falls, then the $(k+1)$ th domino will fall

$\therefore \forall x P(x)$

- All dominoes will fall!



All of your proofs should have the same overall structure

$P(x) \equiv$ *Define the property that you are trying to prove*

Base case: *Prove the “first step onto the ladder.” Typically, but not always, this means proving $P(0)$ or $P(1)$.*

Inductive Hypothesis: *Assume that $P(k)$ is true for an arbitrary k*

Inductive step: *Show that $P(k) \rightarrow P(k + 1)$. That is, prove that once you’re on one step, you can get to the next step. This is where many proofs will differ from one another.*

Conclusion: *Since you’ve proven the base case and $P(k) \rightarrow P(k + 1)$, the claim is true! \square*

Prove that $\sum_{j=1}^n j = \frac{n(n+1)}{2}$

$P(n) \equiv$	$\sum_{j=1}^n j = \frac{n(n+1)}{2}$
Base case:	$P(1): 1(1+1)/2 = 1$ ✓
I.H.:	Assume that $P(k)$ holds for an arbitrary positive integer k
Inductive step:	We will now show that $P(k) \rightarrow P(k+1)$ <ul style="list-style-type: none"> • $1+2+\dots+k = k(k+1)/2$ by I.H. • $1+2+\dots+k+(k+1) = k(k+1)/2 + (k+1)$ $k+1$ to both sides • $1+2+\dots+k+(k+1) = k(k+1)/2 + 2(k+1)/2$ • $1+2+\dots+k+(k+1) = (k^2 + 3k + 2)/2$ • $1+2+\dots+k+(k+1) = (k+1)(k+2)/2$ factoring
Conclusion:	Since we have proved the base case and the inductive case, $\forall n \in \mathbb{Z}^+(P(n))$ by mathematical induction ◻

Induction cannot give us a formula to prove, but can allow us to verify conjectures

Mathematical induction is **not** a tool for discovering new theorems, but rather a powerful way to prove them

Example: Make a conjecture about the sum of the first n odd positive numbers, then prove it.

- $1 = 1$
- $1 + 3 = 4$
- $1 + 3 + 5 = 9$
- $1 + 3 + 5 + 7 = 16$
- $1 + 3 + 5 + 7 + 9 = 25$

The sequence 1, 4, 9, 16, 25, ... appears to be the sequence $\{n^2\}$

Conjecture: The sum of the first n odd positive integers is n^2

Prove that the sum of the first n positive odd integers is n^2

$P(n) \equiv$ The sum of the first n positive odd numbers is n^2

Base case: $P(1): 1 = 1$ ✓

I.H.: Assume that $P(k)$ holds for an arbitrary positive integer k

Inductive step: We will now show that $P(k) \rightarrow P(k+1)$

- $1+3+\dots+(2k-1) = k^2$ by I.H.
- $1+3+\dots+(2k-1)+(2k+1) = k^2+2k+1$ $2k+1$ to both sides
- $1+3+\dots+(2k-1)+(2k+1) = (k+1)^2$ factoring

Note: The k th odd integer is $2k-1$, the $(k+1)$ th odd integer is $2k+1$

Conclusion: Since we have proved the base case and the inductive case, $\forall n \in \mathbb{Z}^+ (P(n))$ by mathematical induction ◻

Prove that the sum $1 + 2 + 2^2 + \dots + 2^n = 2^{n+1} - 1$ for all nonnegative integers n

$$P(n) \equiv \sum_{i=0}^n 2^i = 2^{n+1} - 1$$

Base case: $P(0): 2^0 = 1$ ✓

I.H.: Assume that $P(k)$ holds for an arbitrary **natural number** k

Inductive step: We will now show that $P(k) \rightarrow P(k+1)$

- $1+2+\dots+2^k = 2^{k+1}-1$ by I.H.
- $1+2+\dots+2^k+2^{k+1} = 2^{k+1}-1+2^{k+1}$ 2^{k+1} to both sides
- $1+2+\dots+2^k+2^{k+1} = 2^{k+1}+2^{k+1}-1$ associative law
- $1+2+\dots+2^k+2^{k+1} = 2 \times 2^{k+1}-1$ def'n of ×
- $1+2+\dots+2^k+2^{k+1} = 2^{k+2}-1$ def'n of exp.

Conclusion: Since we have proved the base case and the inductive case, $\forall n \in \mathbf{N}(P(n))$ by mathematical induction ◻

Why does mathematical induction work?

This follows from the **well ordering** axiom

- i.e., Every set of positive integers has a least element

We can prove that mathematical induction is valid using a proof by contradiction.

- Assume that $P(1)$ holds and $P(k) \rightarrow P(k+1)$, but $\neg \forall x P(x)$
- This means that the set $S = \{x \mid \neg P(x)\}$ is nonempty
- By well ordering, S has a least element m with $\neg P(m)$
- Since m is the **least element** of S , $P(m-1)$ is true
- By $P(k) \rightarrow P(k+1)$, $P(m-1) \rightarrow P(m)$
- Since we have $P(m) \wedge \neg P(m)$ this is a **contradiction!**

Result: Mathematical induction is a valid proof method

In-class exercises

Problem 1: Prove that $\sum_{j=0}^n ar^j = \frac{ar^{n+1} - a}{r - 1}$ if $r \neq 1$

Problem 2: Prove that $\sum_{j=1}^n (3j - 2) = \frac{n(3n - 1)}{2}$

Hint: Be sure to

1. Define $P(x)$
2. Prove the base case
3. Make an inductive hypothesis
4. Carry out the inductive step
5. Draw the final conclusion

Prove the formula for the sum of the first n positive squares

$$P(n) \equiv \sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$

Base case: $P(1): 1^2 = \frac{1(1+1)(2+1)}{6}$ ✓

I.H.: Assume that $P(k)$ holds for an arbitrary positive integer k

Inductive step: We will now show that $P(k) \rightarrow P(k+1)$

- $1+4+9+\dots+k^2 = k(k+1)(2k+1)/6$ by I.H.
- $1+4+9+\dots+(k+1)^2 = k(k+1)(2k+1)/6 + (k+1)^2$ (k+1)² to both sides
- $= k(k+1)(2k+1)/6 + 6(k+1)^2/6$ common denom.
- $= (k+1)(2k^2+k+6k+6)/6 = (k+1)(2k^2+7k+6)/6$ factor k+1, mult.
- $= (k+1)(k+2)(2k+3)/6$ factor
- $= (k+1)((k+1)+1)(2(k+1)+1)/6, \therefore P(k+1)$ proved for $k+1$

Conclusion: Since we have proved the base case and the inductive case,
 $\forall n \in \mathbf{Z}^+ (P(n))$ by mathematical induction ◻

Induction can also be used to prove properties other than summations!

$$\leq$$

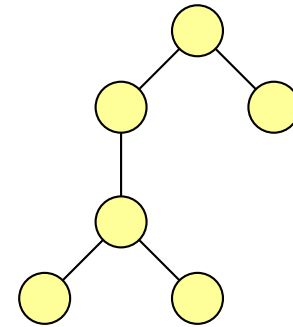
Inequalities

$$\equiv \varphi(p)$$

Divisibility and results from
number theory

$$\in \subseteq$$

Set theory



Algorithms and data structures

Prove that $2^n < n!$ for every positive integer $n \geq 4$

Prelude: The expression $n!$ is called the factorial of n .

Definition: $n! = n \times (n-1) \times \dots \times 3 \times 2 \times 1$

Examples:

- $4! = 4 \times 3 \times 2 \times 1 = 24$
- $5! = 5 \times 4 \times 3 \times 2 \times 1 = 120$
- $6! = 6 \times 5 \times 4 \times 3 \times 2 \times 1 = 720$
- $7! = 7 \times 6 \times 5 \times 4 \times 3 \times 2 \times 1 = 5,040$
- $8! = 8 \times 7 \times 6 \times 5 \times 4 \times 3 \times 2 \times 1 = 40,320$

Note how quickly the factorial of n "grows"



Prove that $2^n < n!$ for every positive integer $n \geq 4$

$P(n) \equiv 2^n < n!$

Base case: $P(4): 2^4 < 4!$ ✓

I.H.: Assume that $P(k)$ holds for an arbitrary integer $k \geq 4$

Inductive step: We will now show that $P(k) \rightarrow P(k+1)$

- $2^k < k!$ by I.H.
- $2 \times 2^k < 2 \times k!$ multiply by 2
- $2^{k+1} < 2 \times k!$ def'n of exp.
- $2^{k+1} < (k+1) \times k!$ since $2 < (k+1)$
- $2^{k+1} < (k+1)!$ def'n of factorial

Conclusion: Since we have proved the base case and the inductive case,
 $\forall n \geq 4 (P(n))$ by mathematical induction ◻

Prove that $n^3 - n$ is divisible by 3 whenever n is a positive integer

$$P(n) \equiv 3 \mid (n^3 - n)$$

Base case: $P(1): 3 \mid 0$ ✓

I.H.: Assume that $P(k)$ holds for an arbitrary positive integer k

Inductive step: We will now show that $P(k) \rightarrow P(k+1)$

- $(k+1)^3 - (k+1) = k^3 + 3k^2 + 3k + 1 - (k+1)$
- $= k^3 + 3k^2 + 2k$
- $= (k^3 - k) + (3k^2 + 3k)$
- $= (k^3 - k) + 3(k^2 + k)$
- Note that $3 \mid (k^3 - k)$ by the I.H. and $3 \mid 3(k^2 + k)$ by definition, so $3 \mid [(k+1)^3 - (k+1)]$

Conclusion: Since we have proved the base case and the inductive case, $\forall n \in \mathbb{Z}^+(P(n))$ by mathematical induction ◻

In-class exercises

Problem 3: Prove that $n^3 + 2n$ is divisible by 3 for any positive integer n

Problem 4: Prove that $6^n - 1$ is divisible by 5 for any positive integer n

Hint: Be sure to

1. Define $P(x)$
2. Prove the base case
3. Make an inductive hypothesis
4. Carry out the inductive step
5. Draw the final conclusion

Prove that if S is a finite set with n elements, then S has 2^n subsets.

$P(n) \equiv$ Set S with cardinality n has 2^n subsets

Base case: $P(0)$: \emptyset has $2^0 = 1$ subsets (i.e., $\emptyset \subseteq \emptyset$) ✓

I.H.: Assume that $P(k)$ holds for an arbitrary natural number k

Inductive step: We will now show that $P(k) \rightarrow P(k+1)$

- Let S be a set of size k
- Assume without loss of generality that $x \notin S$
- Let $T = S \cup \{x\}$, so $|T| = k+1$
- $\forall s \subseteq S$ ($s \subseteq T$) since T is a superset of S
- Furthermore, $\forall s \subseteq S$ ($s \cup \{x\} \subseteq T$) since $x \in T$
- Since S has 2^k subsets by the I.H., T has $2 \times 2^k = 2^{k+1}$ subsets

Conclusion: Since we have proved the base case and the inductive case, $\forall n \in \mathbb{N}(P(n))$ by mathematical induction ◻

Final Thoughts

- Mathematical induction lets us prove universally quantified statements using this inference rule:

$$\begin{array}{l} P(0) \\ \text{For arb. } k, P(k) \rightarrow P(k+1) \\ \hline \therefore \forall x \in \mathbf{N} P(x) \end{array}$$

- Induction is useful for proving:
 - Summations
 - Inequalities
 - Claims about countable sets
 - Theorems from number theory
 - ...
- Next time: Strong induction and recursive definitions (Sections 5.2 & 5.3)