

CS 441: Predicates and Quantifiers

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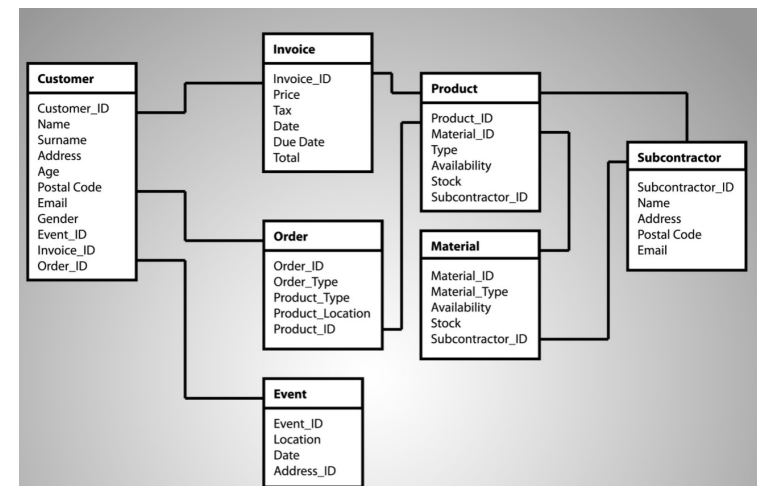


[Key CS Link] Quantifiers

E-commerce Website Developer

Imagine you're the lead developer for an e-commerce website. You need to write a program that checks the inventory and shipping status for customer orders. Your task is to confirm if the following statement is true:

"For every customer, there is a product they have ordered that is currently in stock."



Source: Gemini

Today's topics

- Predicates
- Quantifiers
- Logical equivalences in predicate logic
- Translations using quantifiers



Propositional logic is simple, therefore limited

Propositional logic cannot represent some classes of natural language statements...

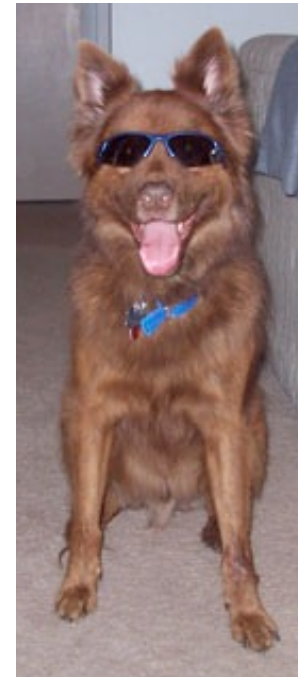


Given: All of my dogs like peanut butter



Propositional logic gives us **no way** to draw the (obvious) conclusion that Kody likes peanut butter!

Given: Kody is one of my dogs



Propositional logic also limits the mathematical truths that we can express and reason about

Consider the following:

- $p_1 \equiv 2$ has no divisors other than 1 and itself
- $p_2 \equiv 3$ has no divisors other than 1 and itself
- $p_3 \equiv 5$ has no divisors other than 1 and itself
- $p_4 \equiv 7$ has no divisors other than 1 and itself
- $p_5 \equiv 11$ has no divisors other than 1 and itself
- ...

This is an inefficient way to reason about the properties of prime numbers!

General problem: Propositional logic has no way of reasoning about instances of general statements.

Historical Context

The previous examples are called **sylogisms**

Aristotle used syllogisms in his *Prior Analytics* to deductively infer new facts from existing knowledge

Major premise

→ All men are mortal

Socrates is a man

← **Minor premise**

∴ Socrates is mortal

Conclusion



Predicate logic allows us to reason about the properties of individual objects and classes of objects

Predicate logic allows us to use propositional functions during our logical reasoning

The diagram shows the equation $P(x) \equiv x^3 > 0$. A large curly brace above the equation spans from $P(x)$ to $x^3 > 0$. A smaller curly brace below $P(x)$ has a line pointing to the word "variable". Another curly brace below $x^3 > 0$ has a line pointing to the word "predicate".

$$P(x) \equiv x^3 > 0$$

variable predicate

Note: A propositional function $P(x)$ has no truth value unless it is evaluated for a given x or set of x s.

Examples

Assume $P(x) \equiv x^3 > 0$. What are the truth values of the following expressions:

- $P(0)$
- $P(23)$
- $P(-42)$

We can express the prime number property using predicate logic:

Predicates can also be defined on more than one variable

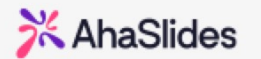
Let $P(x, y) \equiv x + y = 42$. What are the truth values of the following expressions:

- $P(45, -3)$
- $P(23, 23)$
- $P(1, 119)$

Let $S(x, y, z) \equiv x + y = z$. What are the truth values of the following expressions:

- $S(1, 1, 2)$
- $S(23, 24, 42)$
- $S(-9, 18, 9)$

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Quiz question 2 of 5

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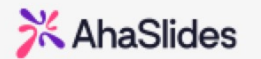
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Quiz question 3 of 5

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Predicates play a central role in program control flow and debugging

If/then statements:

- **if** $x > 17$ ~~then~~ $y = 13$

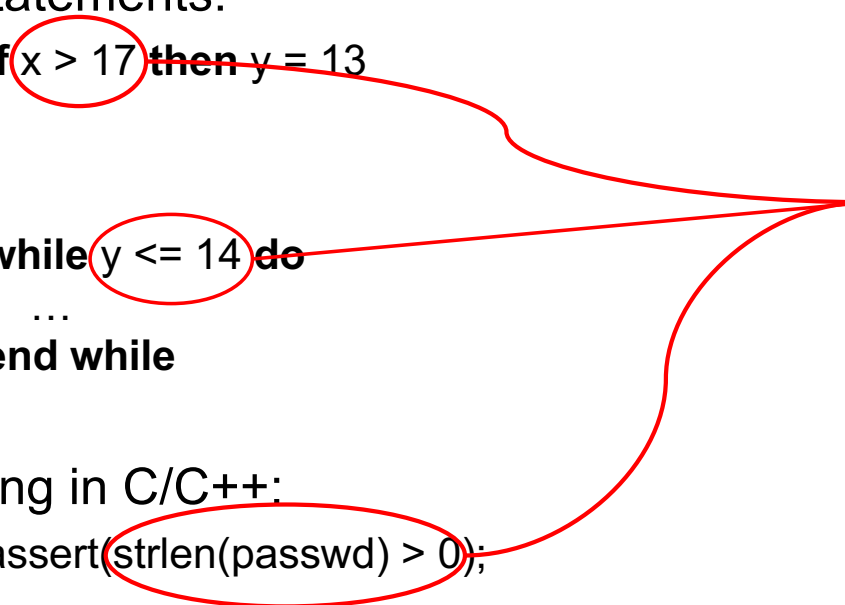
Loops:

- **while** $y \leq 14$ **do**
...
end while

Debugging in C/C++:

- **assert**($\text{strlen}(\text{passwd}) > 0$);

This is a predicate!



Quantifiers allow us to make general statements that turn propositional functions into propositions

In English, we use quantifiers on a regular basis:

- All students can ride the bus for free
- Many people like chocolate
- I enjoy some types of tea
- At least one person will sleep through their final exam

Quantifiers require us to define a **universe of discourse** (also called a **domain**) in order for the quantification to make sense

- “Many like chocolate” doesn’t make sense!

What are the universes of discourse for the above statements?

Universal quantification allows us to make statements about the entire universe of discourse

Examples:

- All of my dogs like peanut butter
- Every even integer is a multiple of two
- For each positive integer x , $2x > x$

Given a propositional function $P(x)$, we express the universal quantification of $P(x)$ as $\forall x P(x)$

What is the truth value of $\forall x P(x)$?

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Examples

All rational numbers are greater than 42

If a natural number is prime, it has no divisors other than 1 and itself

Existential quantifiers allow us to make statements about some objects

Examples:

- **Some** elephants are scared of mice
- **There exist** integers a , b , and c such that the equality $a^2 + b^2 = c^2$ is true
- **There is at least one** person who did better than John on the midterm

Given a propositional function $P(x)$, we express the existential quantification of $P(x)$ as $\exists x P(x)$

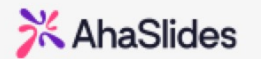
What is the truth value of $\exists x P(x)$?

Examples

The inequality $x + 1 < x$ holds for at least one integer

For some integers, the equality $a^2 + b^2 = c^2$ is True

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Quiz question 1 of 5

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A common idiom in logic: Restricting the domain of quantification

The square of every natural number less than 4 is no more than 9

- *Domain:* natural numbers
- *Statement:* $\forall x < 4 (x^2 \leq 9)$
- *Truth value:* True

This is equivalent to writing
 $\forall x [(x < 4) \rightarrow (x^2 \leq 9)]$



Some integers between 0 and 6 are prime

- *Domain:* Integers
- *Propositional function:* $P(x) \equiv$ "x is prime"
- *Statement:* $\exists 0 \leq x \leq 6 P(x)$
- *Truth value:* True

This is equivalent to writing
 $\exists x [(0 \leq x \leq 6) \wedge P(x)]$



Precedence of quantifiers

The universal and existential quantifiers have the **highest precedence** of all logical operators

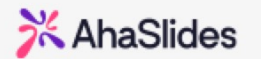
For example:

- $\forall x P(x) \rightarrow Q(x)$ actually means $(\forall x P(x)) \rightarrow Q(x)$
- $\exists x P(x) \wedge Q(x)$ actually means $(\exists x P(x)) \wedge Q(x)$ 

x is undefined outside!

When needed, use parentheses to disambiguate a quantifier's scope

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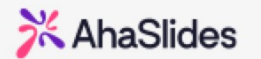
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We can extend the notion of logical equivalence to expressions containing predicates or quantifiers

Definition: Two statements involving predicates and quantifiers are **logically equivalent** iff they take on the same truth value *regardless* of which predicates are substituted into these statements and which domains of discourse are used.

Prove: $\exists x [P(x) \vee Q(x)] \equiv \exists x P(x) \vee \exists x Q(x)$

We must prove each “direction” of the equivalence. Assume that P and Q have the same domain.

First, prove $\exists x [P(x) \vee Q(x)] \rightarrow \exists x P(x) \vee \exists x Q(x)$:

- If $\exists x [P(x) \vee Q(x)]$ is **True**, this means that there is some value v in the domain such that either $P(v)$ is **True** or $Q(v)$ is **True**
- If $P(v)$ is **True**, then $\exists x P(x)$ is **True** and $[\exists x P(x) \vee \exists x Q(x)]$ is **True**
- If $Q(v)$ is **True**, then $\exists x Q(x)$ is **True** and $[\exists x P(x) \vee \exists x Q(x)]$ is **True**
- Thus $\exists x [P(x) \vee Q(x)] \rightarrow \exists x P(x) \vee \exists x Q(x)$

Prove: $\exists x [P(x) \vee Q(x)] \equiv \exists x P(x) \vee \exists x Q(x)$

Then, prove $\exists x P(x) \vee \exists x Q(x) \rightarrow \exists x [P(x) \vee Q(x)]$:

- If $\exists x P(x) \vee \exists x Q(x)$ is **True**, this means that there is some value v in the domain such that either $P(v)$ is **True** or $Q(v)$ is **True**
- If $P(v)$ is **True**, then $\exists x [P(x) \vee Q(x)]$ is **True**
- If $Q(v)$ is **True**, then $\exists x [P(x) \vee Q(x)]$ is **True**
- Thus $\exists x P(x) \vee \exists x Q(x) \rightarrow \exists x [P(x) \vee Q(x)]$

Since $\exists x [P(x) \vee Q(x)] \rightarrow \exists x P(x) \vee \exists x Q(x)$ **and**
 $\exists x P(x) \vee \exists x Q(x) \rightarrow \exists x [P(x) \vee Q(x)]$ **then**
 $\exists x [P(x) \vee Q(x)] \equiv \exists x P(x) \vee \exists x Q(x)$.

We also have DeMorgan's laws for quantifiers

Negation over universal quantifier: $\neg(\forall x P(x)) \equiv \exists x (\neg P(x))$

Negation over existential quantifier: $\neg(\exists x P(x)) \equiv \forall x (\neg P(x))$

These are **very** useful logical equivalences, so let's prove one of them...

Prove: $\neg \forall x P(x) \equiv \exists x \neg P(x)$

- $\neg \forall x P(x) \rightarrow \exists x \neg P(x)$
 - $\neg \forall x P(x)$ is **True** if and only if $\forall x P(x)$ is **False**
 - $\forall x P(x)$ is **False** if and only if there is some v such that $P(v)$ is **False**
 - $P(v)$ is **False** is equivalent to $\neg P(v)$ is **True**
 - If $\neg P(v)$ is **True**, then $\exists x \neg P(x)$
- $\exists x \neg P(x) \rightarrow \neg \forall x P(x)$
 - $\exists x \neg P(x)$ is **True** if and only if there is some v such that $\neg P(v)$ is **True**
 - If $\neg P(v)$ is **True**, then clearly $P(x)$ does not hold for all possible values in the domain and thus we have $\neg \forall x P(x)$

Therefore $\neg \forall x P(x) \equiv \exists x \neg P(x)$.

Translations from English

To translate English sentences into logical expressions:

1. Rewrite the sentence to make it easier to translate
2. Determine the appropriate quantifiers to use
3. Look for words that indicate logical operators
4. Formalize sentence fragments
5. Put it all together

Example: At least one person in this classroom is named Nils and has lived in Pittsburgh for 12 years

② *Existential quantifier*

① **Rewrite:** There exists at least one person who is in this classroom, is named Nils, and has lived in Pittsburgh for 12 years

③ *Conjunctions*

④ **Formalize:**

- $C(x) \equiv$ "x is in this classroom"
- $N(x) \equiv$ "x is named Nils"
- $P(x) \equiv$ "x has lived in Pittsburgh for 12 years"

⑤ **Final expression:** $\exists x [C(x) \wedge N(x) \wedge P(x)]$

Example: If a student is taking CS441, then they have taken high school algebra

②

Universal quantifier

① **Rewrite:** For all students, if a student is in CS 441, then they have taken high school algebra

④ **Formalize:**

- $C(x) \equiv$ "x is taking CS441"
- $H(x) \equiv$ "x has taken high school algebra"

③

Implication

⑤ **Final expression:** $\forall x [C(x) \rightarrow H(x)]$

Negate the previous example

$$\neg \forall x [C(x) \rightarrow H(x)] \equiv$$

Translate back into English:

- There is a student taking CS441 that has not taken high school algebra!

Example: Jane enjoys drinking some types of tea

Rewrite: There exist some types of tea that Jane enjoys drinking

Formalize:

- $T(x) \equiv$ “x is a type of tea”
- $D(x) \equiv$ “Jane enjoys drinking x”

Final expression: $\exists x [T(x) \wedge D(x)]$

Negate the previous example:

$$\neg \exists x [T(x) \wedge D(x)]$$

In-class Activities

Activity 1: Translate the following sentences into logical expressions. Remember to state all domains [[miro](#)].

- a) Some cows have black spots
- b) At least one student likes to watch football or ice hockey

Activity 2: Negate the translated expressions from Activity 1. Translate these back into English. [[miro](#)]



Steps:

1. Introduce to a classmate
2. Work in pairs on the exercise
3. Submit answers on miro
4. Volunteers to share answers

Final Thoughts

- The simplicity of propositional logic makes it unsuitable for solving certain types of problems
- Predicate logic makes use of
 - Propositional functions to describe properties of objects
 - The universal quantifier to assert properties of **all** objects within a given domain
 - The existential quantifier to assert properties of **some** objects within a given domain
- Predicate logic can be used to reason about relationships between objects and classes of objects
- **Next lecture:**
 - Applications of predicate logic and nested quantifiers
 - Please read section 1.5